

CS244a: An Introduction to Computer Networks

Handout 8: Congestion Avoidance and Active Queue Management



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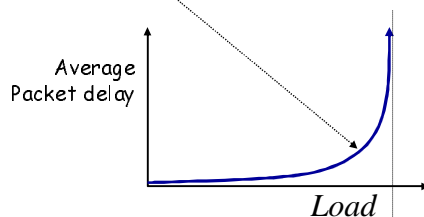
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Congestion Avoidance

- ❖ TCP reacts to congestion *after* it takes place. The data rate changes rapidly and the system is barely stable (or is even unstable).
- ❖ Can we *predict* when congestion is about to happen and avoid it? E.g. by detecting the knee of the curve.



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Congestion Avoidance Schemes

❖ Router-based Congestion Avoidance:

❖ DECbit:

- Routers explicitly notify sources about congestion.

❖ Random Early Detection (RED):

- Routers implicitly notify sources by dropping packets.
- RED drops packets at random, and as a function of the level of congestion.

❖ Host-based Congestion Avoidance

- ❖ Source monitors changes in RTT to detect onset of congestion.
- ❖ A variety of algorithms are described in Section 6.4.3.

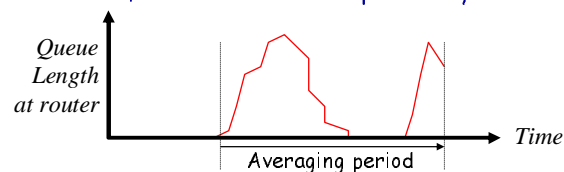
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DECbit

- ❖ Each packet has a "Congestion Notification" bit called the DECbit in its header.
- ❖ If any router on the path is congested, it sets the DECbit.
 - ❖ Set if average queue length ≥ 1 packet, averaged since the start of the previous busy cycle.
- ❖ To notify the source, the destination copies DECbit into ACK packets.
- ❖ Source adjusts rate to avoid congestion.
 - ❖ Counts fraction of DECbits set in each window.
 - ❖ If $< 50\%$ set, increase rate additively.
 - ❖ If $\geq 50\%$ set, decrease rate multiplicatively.



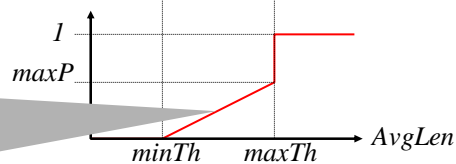
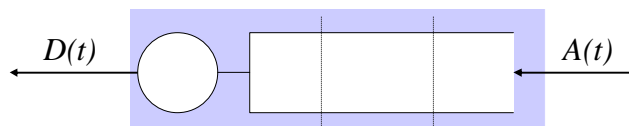
Random Early Detection (RED)

- ❖ RED is based on DECbit, and was designed to work well with TCP.
- ❖ RED implicitly notifies sender by dropping packets.
- ❖ Drop probability is increased as the *average* queue length increases.
- ❖ (Geometric) moving average of the queue length is used so as to detect long term congestion, yet allow short term bursts to arrive.

$$AvgLen_{n+1} = (1-\alpha) \times AvgLen_n + \alpha \times Length_n$$

$$\text{i.e. } AvgLen_{n+1} = \sum_{i=1}^n Length_i (\alpha)(1-\alpha)^{n-i}$$

RED Drop Probabilities



If $minTh < AvgLen < maxTh$:

$$\hat{p}_{AvgLen} = maxP \left\{ \frac{AvgLen - minTh}{maxTh - minTh} \right\}$$

$$Pr(\text{Drop Packet}) = \frac{\hat{p}_{AvgLen}}{1 - count \times \hat{p}_{AvgLen}}$$

count counts how long we've been in $minTh < AvgLen < maxTh$ since we last dropped a packet. i.e. drops are spaced out in time, reducing likelihood of re-entering slow-start.

Properties of RED

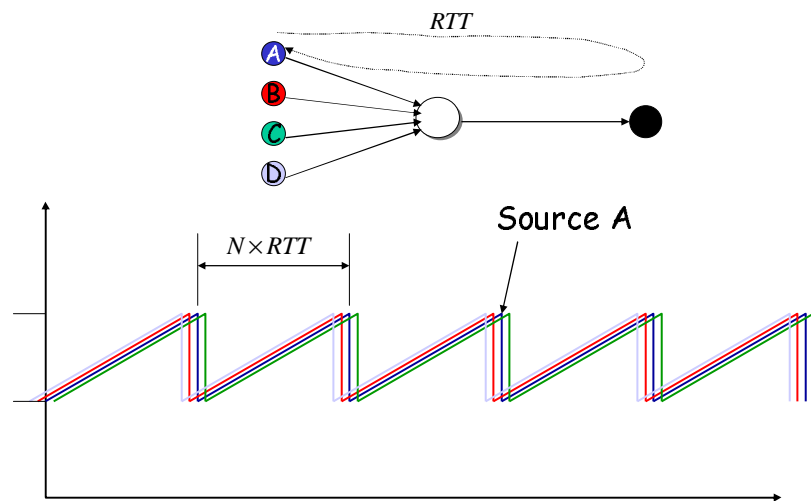
- ❖ Drops packets before queue is full, in the hope of reducing the rates of some flows.
- ❖ Drops packet for each flow *roughly* in proportion to its rate.
- ❖ Drops are spaced out in time.
- ❖ Because it uses average queue length, RED is tolerant of bursts.
- ❖ Random drops hopefully desynchronize TCP sources.

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Synchronization of sources

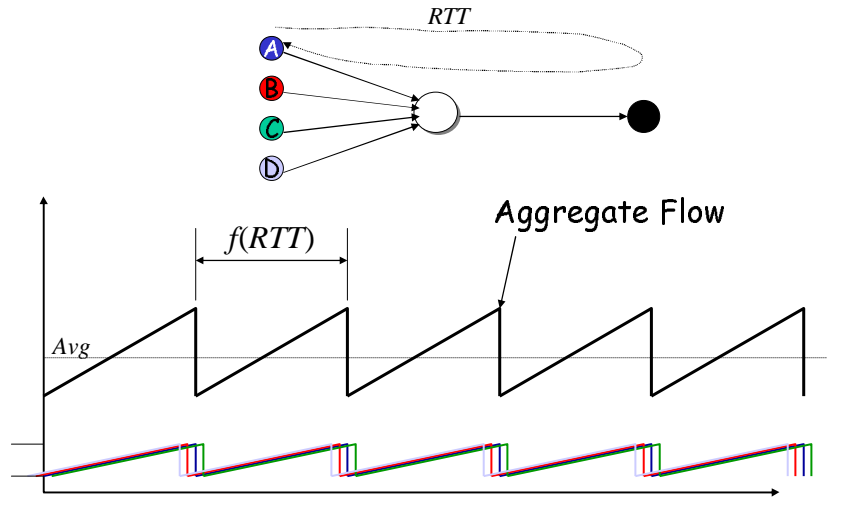


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Synchronization of sources

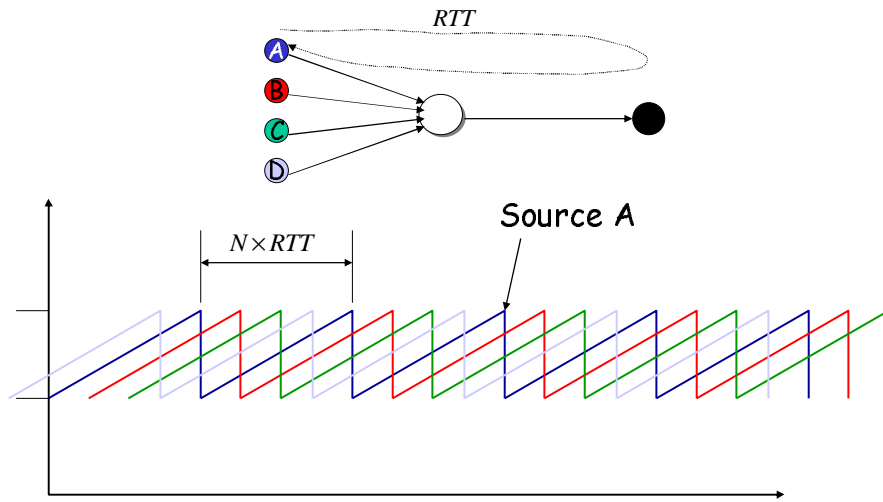


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Desynchronized sources



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Desynchronized sources

